Scheduling the CPU is clearly a difficult problem

Scheduling the CPU is clearly a difficult problem

It requires the collection and manipulation of many statistics about processes

Scheduling the CPU is clearly a difficult problem

It requires the collection and manipulation of many statistics about processes

Scheduling one resource (the CPU) is hard enough

Scheduling the CPU is clearly a difficult problem

It requires the collection and manipulation of many statistics about processes

Scheduling one resource (the CPU) is hard enough

We now look at a new problem that arises when we want to schedule *multiple* resources



Processes compete for resources like disks and network and the OS mediates this



Processes compete for resources like disks and network and the OS mediates this

To read from a disk, a process must call the OS kernel and wait for the kernel to reply

## Terminology

When we say "a process waits for the kernel" we mean, of course, something entirely different

# Terminology

When we say "a process waits for the kernel" we mean, of course, something entirely different

What actually happens is the kernel marks the process as blocked, and does not consider it for scheduling until the requested resource has arrived

# Terminology

When we say "a process waits for the kernel" we mean, of course, something entirely different

What actually happens is the kernel marks the process as blocked, and does not consider it for scheduling until the requested resource has arrived

There is no "waiting" happening: the process does not run when blocked



Processes compete for resources like disks and network and the OS mediates this

To read from a disk, a process must call the OS kernel and wait for the kernel to reply

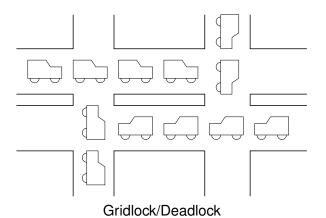


Processes compete for resources like disks and network and the OS mediates this

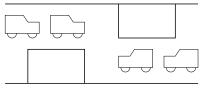
To read from a disk, a process must call the OS kernel and wait for the kernel to reply

Sometimes the delay is infinite!









Gridlock/Deadlock



This can happen in an OS

This can happen in an OS

Process  $P_1$  wants to copy some data from disk  $D_1$  to disk  $D_2$ , while process  $P_2$  wants to copy some data from disk  $D_2$  to disk  $D_1$ 

This can happen in an OS

Process  $P_1$  wants to copy some data from disk  $D_1$  to disk  $D_2$ , while process  $P_2$  wants to copy some data from disk  $D_2$  to disk  $D_1$ 

Initially P<sub>1</sub> is running and makes a request for access to D<sub>2</sub>

This can happen in an OS

Process  $P_1$  wants to copy some data from disk  $D_1$  to disk  $D_2$ , while process  $P_2$  wants to copy some data from disk  $D_2$  to disk  $D_1$ 

- Initially P<sub>1</sub> is running and makes a request for access to D<sub>2</sub>
- The OS takes over and gives *P*<sub>1</sub> exclusive access to *D*<sub>2</sub> (not a smart OS)

This can happen in an OS

Process  $P_1$  wants to copy some data from disk  $D_1$  to disk  $D_2$ , while process  $P_2$  wants to copy some data from disk  $D_2$  to disk  $D_1$ 

- Initially P<sub>1</sub> is running and makes a request for access to D<sub>2</sub>
- The OS takes over and gives *P*<sub>1</sub> exclusive access to *D*<sub>2</sub> (not a smart OS)
- The OS decides to run P2

This can happen in an OS

Process  $P_1$  wants to copy some data from disk  $D_1$  to disk  $D_2$ , while process  $P_2$  wants to copy some data from disk  $D_2$  to disk  $D_1$ 

- Initially P<sub>1</sub> is running and makes a request for access to D<sub>2</sub>
- The OS takes over and gives *P*<sub>1</sub> exclusive access to *D*<sub>2</sub> (not a smart OS)
- The OS decides to run P2
- P<sub>2</sub> runs and makes a request for access to D<sub>1</sub>

This can happen in an OS

Process  $P_1$  wants to copy some data from disk  $D_1$  to disk  $D_2$ , while process  $P_2$  wants to copy some data from disk  $D_2$  to disk  $D_1$ 

- Initially P<sub>1</sub> is running and makes a request for access to D<sub>2</sub>
- The OS takes over and gives *P*<sub>1</sub> exclusive access to *D*<sub>2</sub> (not a smart OS)
- The OS decides to run P2
- P<sub>2</sub> runs and makes a request for access to D<sub>1</sub>
- The OS takes over and gives P<sub>2</sub> exclusive access to D<sub>1</sub>

• The OS decides to run P<sub>1</sub>

- The OS decides to run P1
- P<sub>1</sub> runs and makes a request for access to D<sub>1</sub>

- The OS decides to run P<sub>1</sub>
- P<sub>1</sub> runs and makes a request for access to D<sub>1</sub>
- The OS takes over and notices P<sub>2</sub> has locked D<sub>1</sub>, so P<sub>1</sub> must wait until P<sub>2</sub> has finished with it; P<sub>1</sub> moves to state blocked

- The OS decides to run P<sub>1</sub>
- P<sub>1</sub> runs and makes a request for access to D<sub>1</sub>
- The OS takes over and notices P<sub>2</sub> has locked D<sub>1</sub>, so P<sub>1</sub> must wait until P<sub>2</sub> has finished with it; P<sub>1</sub> moves to state blocked
- The OS decides to run  $P_2$ : it can't run  $P_1$  as it is blocked

- The OS decides to run P<sub>1</sub>
- P<sub>1</sub> runs and makes a request for access to D<sub>1</sub>
- The OS takes over and notices P<sub>2</sub> has locked D<sub>1</sub>, so P<sub>1</sub> must wait until P<sub>2</sub> has finished with it; P<sub>1</sub> moves to state blocked
- The OS decides to run P<sub>2</sub>: it can't run P<sub>1</sub> as it is blocked
- P<sub>2</sub> runs and makes a request for access to D<sub>2</sub>

- The OS decides to run P<sub>1</sub>
- P<sub>1</sub> runs and makes a request for access to D<sub>1</sub>
- The OS takes over and notices P<sub>2</sub> has locked D<sub>1</sub>, so P<sub>1</sub> must wait until P<sub>2</sub> has finished with it; P<sub>1</sub> moves to state blocked
- The OS decides to run P<sub>2</sub>: it can't run P<sub>1</sub> as it is blocked
- P<sub>2</sub> runs and makes a request for access to D<sub>2</sub>
- The OS takes over and notices P<sub>1</sub> has locked D<sub>2</sub>, so P<sub>2</sub> must wait until P<sub>1</sub> has finished with it; P<sub>2</sub> moves to state blocked

- The OS decides to run P<sub>1</sub>
- P<sub>1</sub> runs and makes a request for access to D<sub>1</sub>
- The OS takes over and notices P<sub>2</sub> has locked D<sub>1</sub>, so P<sub>1</sub> must wait until P<sub>2</sub> has finished with it; P<sub>1</sub> moves to state blocked
- The OS decides to run  $P_2$ : it can't run  $P_1$  as it is blocked
- P2 runs and makes a request for access to D2
- The OS takes over and notices P<sub>1</sub> has locked D<sub>2</sub>, so P<sub>2</sub> must wait until P<sub>1</sub> has finished with it; P<sub>2</sub> moves to state blocked
- Now both *P*<sub>1</sub> and *P*<sub>2</sub> are blocked and the OS can't run either process!

 $P_1$  can't run until  $D_1$  is free, but  $D_1$  won't be free until  $P_2$  runs  $P_2$  can't run until  $D_2$  is free, but  $D_2$  won't be free until  $P_1$  runs

 $P_1$  can't run until  $D_1$  is free, but  $D_1$  won't be free until  $P_2$  runs  $P_2$  can't run until  $D_2$  is free, but  $D_2$  won't be free until  $P_1$  runs This is called *deadlock* 

 $P_1$  can't run until  $D_1$  is free, but  $D_1$  won't be free until  $P_2$  runs  $P_2$  can't run until  $D_2$  is free, but  $D_2$  won't be free until  $P_1$  runs This is called *deadlock* 

Deadlock can happen on any kind of shared resources that require exclusive access

 $P_1$  can't run until  $D_1$  is free, but  $D_1$  won't be free until  $P_2$  runs  $P_2$  can't run until  $D_2$  is free, but  $D_2$  won't be free until  $P_1$  runs This is called *deadlock* 

Deadlock can happen on any kind of shared resources that require exclusive access

And with more than two processes: think of three or more processes in a circle



A formal definition:



A formal definition:

A set of processes D is deadlocked if

- 1. each process  $P_i$  in D is blocked on some event  $e_i$
- 2. event  $e_i$  can only be caused by some process in D



Note that you can only get deadlock if

Note that you can only get deadlock if

• there is more than one resource

Note that you can only get deadlock if

- there is more than one resource
- there is more than one process

Note that you can only get deadlock if

- there is more than one resource
- there is more than one process<sup>1</sup>

<sup>1</sup> It could technically happen with just one process, but that would be quite dumb programming to request for a resource you already have

Note that you can only get deadlock if

- there is more than one resource
- there is more than one process<sup>12</sup>

<sup>1</sup> It could technically happen with just one process, but that would be quite dumb programming to request for a resource you already have <sup>2</sup> I've seen it happen

< 17 >