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As declarative languages don't specify how to do something, the system is free to do things in the most efficient way it can: and this includes in parallel

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Plus a lot of other less obvious problems



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Which may lure them into a false sense that they understand what they are doing in the parallel version of the language

And the legacy language likely has no in-built prevention of things that are safe sequentially, but dangerous in parallel (e.g., updating a shared variable)

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Rust: designed to be a memory-safe replacement for C



- Occam: You shoot both your feet with several guns at once
- Go: To shoot yourself in the foot you must first import the unsafe package
- Rust: you try to shoot yourself in the foot, but you can't as the gun has immutably borrowed your foot

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And they typically have a lot of parallelism (and other!) bugs



But there are many large, successful parallel systems built using these languages



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But the economics of doing this (time to train programmers, time to write and debug new code, etc.) means it rarely happens



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There is much more to be said. There is a whole final-year Unit on parallelism!

Some languages manage memory allocation and deallocation for you, some don't

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For example, some languages have garbage collection, while others require the programmer to take care over memory deallocation

In code

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bigclass x = new bigclass(1);
x = new bigclass(2);
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(setq	x	(make	<bigclass></bigclass>	2))

They both allocate memory to store an instance of bigclass (initialised with 1); then they allocate more memory to store another instance of bigclass (initialised with 2)

But the memory allocated to the object 1 is no longer accessible to the program as the reference to the object stored in the variable x has been overwritten

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Otherwise, that memory is just garbage. We can use a *garbage collector* to search out such inaccessible memory and reclaim it and reuse it

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Often, the user program has to stop running while the GC does its thing as the GC may move values around in memory to tidy up the free spaces

Though well-designed GC code runs very quickly, or, in a few systems, in parallel with the application code

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Thus it is buggy C++ with a *memory leak*: the memory is never recovered and reused

With MMM the programmer needs to explicitly allocate memory for values and then free (deallocate) that memory when they are done with it

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int *array = malloc(4*sizeof(int)); // allocate
...
... do stuff with array
...
free(array); // deallocate
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Remember, in real code, the allocate and free can be thousands of lines of code apart, and written by different programmers; or there might be more than one place where allocation or deallocation could be done; and so on

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Writing into the array is overwriting the bytes used by some other value in the system

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**Exercise** Find the time it takes to create an object in Java, or your favourite language

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That is to say, the *application programmer* doesn't have to take care

(The programmer who implemented the GC in the runtime definitely needed to care!)

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MMM: allows precise memory management as the programmer will explicitly allocate and deallocate memory, but also facilitates buggy programming as it is easy for the programmer to get it wrong

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MMM: allows precise memory management as the programmer will explicitly allocate and deallocate memory, but also facilitates buggy programming as it is easy for the programmer to get it wrong

The programmer can tune the use of memory for their application

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**Exercise** Is this the best of both worlds or the worst of both worlds?

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"Around 70 percent of all the vulnerabilities in Microsoft products addressed through a security update each year are memory safety issues" Matt Miller, Microsoft security engineer, Feb 2019

Memory safety bugs in C and C++ continue to be the most-difficult-to-address source of incorrectness. We invest a great deal of effort and resources into detecting, fixing, and mitigating this class of bugs, and these efforts are effective in preventing a large number of bugs from making it into Android releases. Yet in spite of these efforts, memory safety bugs continue to be a top contributor of stability issues, and consistently represent approximately 70% of Android's high severity security vulnerabilities.

Jeff Vander Stoep and Stephen Hines, Android Team, April 2021

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- double free
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None of these are applicable to a GC language

Other memory errors include:

- reading uninitialised memory
- accessing beyond the bounds of allocated memory, e.g., beyond the ends of a vector

affecting both GC and MMM

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Some languages have memory access checking to avoid these kinds of errors, but checking will slow the running of your program down

For example, x[n] = 42;

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- Safe, but slower
- Unchecked: the code just does the assignment
- Faster, but allows bugs

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- (And C vectors don't include information about their length, anyway)

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**Exercise** What does x [-1] do in Python?

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Exercise And Java, Rust, C++, ...?

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Some tools work at compile time; some at run time

**Exercise** Read about valgrind, Allinea, AddressSanitizer and other runtime memory checking tools

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**Exercise** Find out whether your favourite language checks or not; if it checks, what are the overheads?