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Note we now have *three* kinds of addresses: physical, network and human





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Exercise The file lives on as /etc/hosts under Unix. Look at this file

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That is, machines running a DNS server program

Aside

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Exercise Find out the services your home Access Point runs





 ${\tt bath.ac.uk}$ is a network in the domain ${\tt ac.uk},$ and so on



bath.ac.uk is a network in the domain ac.uk, and so on

 ${\tt ac.uk}$ is the name for the JANET network, and is in the domain ${\tt uk}$



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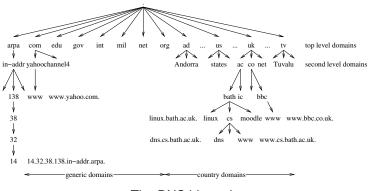
uk is in the domain . (*dot* or *root*)

Each node in this tree is called a *label*



www.llanfairpwllgwyngyllgogerychwyrndrobwllllantysiliogogogoch.com is a valid name

Labels may be up to 63 bytes long



The DNS hierarchy



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But the important bit is that they *delegate* management of lower labels to other organisations

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Labels under cs.bath.ac.uk are managed by the Department of Computer Science DDAT



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So, starting a dot, we can work our way down the tree to find the machine we want

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A few years ago, there was a convention of having an off-site replica, too, e.g., Bath used ns2.ja.net



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If a host on the Bath network needs to look up a local name, linux.bath.ac.uk, say, it sends a request to one of the local servers, e.g., ns2.bath.ac.uk

And the local server will look it up and return the answer



How does that host know which machine to ask? In particular, how does it know the local DNS server's IP address?



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So it can send the DNS request to one of these servers



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If a host in the University requires a name lookup of a non-local name, say news.bbc.co.uk, it sends a DNS request to the **local** DNS server, ns1.bath.ac.uk, say



In this example, the Bath server does not have responsibility for the name news.bbc.co.uk, but it will helpfully do a lookup for us. It will run down the DNS tree in a *recursive lookup*



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The Bath server ns1 sends a *start of authority* (SOA) request to a random *top level* server to find who is responsible for the uk label

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This a machine that is responsible for uk labels



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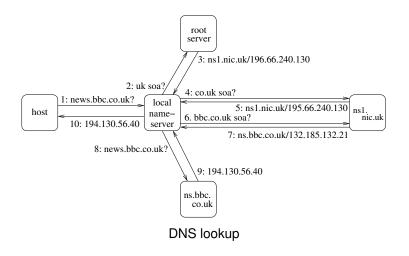
It gets the IP address 194.130.56.40



The server sends an *address* (A) request for news.bbc.co.uk to ns.bbc.co.uk

It gets the IP address 194.130.56.40

Finally, the Bath server can relay this back to the original requesting host





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The next request for news.bbc.co.uk can be answered directly by the local nameserver

Similarly, given a request for www.bbc.co.uk, it can go directly to ns.bbc.co.uk



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Short lived associations will have a short TTL



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This means a faster response and a decrease in network traffic

And less work for the host

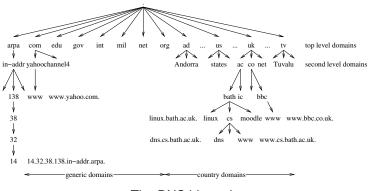


Sometimes we want to do the reverse lookup: given an IP address find a name



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There is a part of the tree dedicated to this with TLD arpa



The DNS hierarchy





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We get clan.bath.ac.uk



DNS does

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And many more: about 50 in total, though few are used frequently



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• One address can have several names

DNS

Both news.bbc.co.uk and newswww.bbc.net.uk resolve to 212.58.226.33

news.bbc.co.uk. 1619 IN CNAME newswww.bbc.net.uk. newswww.bbc.net.uk. 77 IN A 212.58.226.33

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This allows us tricks, like moving the Web server to different hosts, even different ISPs or different countries, but keeping its public name the same



 One name can have several IP addresses associated. This allows *load sharing*, e.g., a Web server can actually be several machines spread about anywhere in the world

www.microsoft.com.	68	IN	CI
toggle.www.ms.akadns.ne	et. 285	IN	CI
g.www.ms.akadns.net.	285	IN	CI
lb1.www.ms.akadns.net.	285	IN	А
lb1.www.ms.akadns.net.	285	IN	А
lb1.www.ms.akadns.net.	285	IN	А
lb1.www.ms.akadns.net.	285	IN	А

IN CNAME toggle.www.ms.akadns.net.
IN CNAME g.www.ms.akadns.net.
IN CNAME lb1.www.ms.akadns.net.
IN A 207.46.19.190
IN A 207.46.19.254
IN A 207.46.192.254
IN A 207.46.193.254

www.microsoft.com.	68	IN	CNAME	<pre>toggle.www.ms.akadns.net.</pre>
toggle.www.ms.akadns.ne	t. 285	IN	CNAME	g.www.ms.akadns.net.
g.www.ms.akadns.net.	285	IN	CNAME	lb1.www.ms.akadns.net.
lb1.www.ms.akadns.net.	285	IN	Α	207.46.19.190
lb1.www.ms.akadns.net.	285	IN	Α	207.46.19.254
lb1.www.ms.akadns.net.	285	IN	Α	207.46.192.254
lb1.www.ms.akadns.net.	285	IN	А	207.46.193.254

www.microsoft.com is an alias: its canonical name is toggle.www.ms.akadns.net;

www.microsoft.com.	68	IN	CNAME	<pre>toggle.www.ms.akadns.net.</pre>
toggle.www.ms.akadns.ne	t. 285	IN	CNAME	g.www.ms.akadns.net.
g.www.ms.akadns.net.	285	IN	CNAME	lb1.www.ms.akadns.net.
lb1.www.ms.akadns.net.	285	IN	Α	207.46.19.190
lb1.www.ms.akadns.net.	285	IN	Α	207.46.19.254
lb1.www.ms.akadns.net.	285	IN	Α	207.46.192.254
lb1.www.ms.akadns.net.	285	IN	А	207.46.193.254

But that is an alias for g.www.ms.akadns.net;

www.microsoft.com.	68	IN	CNAME	<pre>toggle.www.ms.akadns.net.</pre>
toggle.www.ms.akadns.net. 285		IN	CNAME	g.www.ms.akadns.net.
g.www.ms.akadns.net.	285	IN	CNAME	lb1.www.ms.akadns.net.
lb1.www.ms.akadns.net.	285	IN	Α	207.46.19.190
lb1.www.ms.akadns.net.	285	IN	Α	207.46.19.254
lb1.www.ms.akadns.net.	285	IN	Α	207.46.192.254
lb1.www.ms.akadns.net.	285	IN	А	207.46.193.254

And that is an alias for lb1.www.ms.akadns.net;

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toggle.www.ms.akadns.ne	et. 285	IN	CNAME	g.www.ms.akadns.net.
g.www.ms.akadns.net.	285	IN	CNAME	lb1.www.ms.akadns.net.
lb1.www.ms.akadns.net.	285	IN	Α	207.46.19.190
lb1.www.ms.akadns.net.	285	IN	А	207.46.19.254
lb1.www.ms.akadns.net.	285	IN	А	207.46.192.254
lb1.www.ms.akadns.net.	285	IN	Α	207.46.193.254

And that refers to four different addresses



And a DNS server can give out different lookups dependent on who is asking



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A recent lookup for www.microsoft.com returned 104.78.177.250 from one location, and 2.18.85.172 from another



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This allows for load sharing; and also can be used for *geofencing*: giving different services to clients in different place, e.g., videos that are licensed only for certain regions



Exercise Redo these lookups to see how they currently turn out

Exercise Compare having DNS give out multiple IP addresses for a given name against giving out different addresses for different clients against using anycast addresses





Other languages will have similar APIs



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Exercise Old code using the obsolete v4-specific gethostbyname() is one of the sticking points in the transition to IPv6. Read about this



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Under Linux the dig tool can be used to do direct lookups

% dig news.bbc.co.uk :: Got answer: ;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 15281 ;; flags: qr rd ra; QUERY: 1, ANSWER: 2, AUTHORITY: 2, ADDITIONAL: 2 ;; QUESTION SECTION: :news.bbc.co.uk. ΤN Α :: ANSWER SECTION: news.bbc.co.uk. ΤN CNAME newswww.bbc.net.uk. 193 newswww.bbc.net.uk. 76 ΙN Α 212.58.226.73 :: AUTHORITY SECTION:

bbc.net.uk. 85129 ΙN NS ns0.thdo.bbc.co.uk. 85129 NS bbc.net.uk. ΤN ns0.rbsov.bbc.co.uk. ;; ADDITIONAL SECTION: ns0.thdo.bbc.co.uk. 9490 ΤN 212.58.224.20 Α ns0.rbsov.bbc.co.uk. 9490 ΙN Α 212.58.227.47



Quite often a DNS server will reply with more information than we requested, e.g., the lookup of the CNAME newswww.bbc.net.uk in the above



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This means we don't have to do an additional query to get the actual IP address we were looking for

<pre>% dig -x 212.58.226.73 ;; Got answer: ;; ->>HEADER<<- opcode: QUERY, ;; flags: qr rd ra; QUERY: 1, A</pre>		-	
;; QUESTION SECTION: ;73.226.58.212.in-addr.arpa.	IN	PTR	
;; ANSWER SECTION: 73.226.58.212.in-addr.arpa. 500	81 IN	PTR	newslb305.telhc.bbc.co.uk.
;; AUTHORITY SECTION:			
226.58.212.in-addr.arpa. 50081	IN	NS	ns1.thdo.bbc.co.uk.
226.58.212.in-addr.arpa. 50081	IN	NS	ns1.thny.bbc.co.uk.
226.58.212.in-addr.arpa. 50081	IN	NS	ns1.bbc.co.uk.
226.58.212.in-addr.arpa. 50081	IN	NS	ns.bbc.co.uk.
;; ADDITIONAL SECTION:			
ns1.bbc.co.uk. 311	IN	Α	132.185.132.21
ns1.thny.bbc.co.uk. 32106	IN	Α	212.58.227.48
ns1.thdo.bbc.co.uk. 33051	IN	Α	212.58.224.21

% dig +trace www.google.com

	180695	IN	NS	E.ROOT-SERVERS.NET.
	180695	IN	NS	J.ROOT-SERVERS.NET.
	180695	IN	NS	I.ROOT-SERVERS.NET.
com.	172800	IN	NS	A.GTLD-SERVERS.NET.
com.	172800	IN	NS	B.GTLD-SERVERS.NET.
com.	172800	IN	NS	C.GTLD-SERVERS.NET.
google.com.	172800	IN	NS	ns1.google.com.
google.com.	172800	IN	NS	ns2.google.com.
google.com.	172800	IN	NS	ns3.google.com.
www.google.com.	604800	IN	CNAME	www.l.google.com.
l.google.com.	86400	IN	NS	b.l.google.com.
l.google.com.	86400	IN	NS	d.l.google.com.

. . .

DNS labels can be in arbitrary character sets, not just Latin

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A non-Latin DNS name

From blog.icann.org



DNS is a very successful protocol: fast and it's relatively easy for system administrators to manage the DNS servers



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DNS names are a big source of money, so often a source of contention over who should be in control of what



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DNS names are a big source of money, so often a source of contention over who should be in control of what

Exercise Read about the controversies behind the introduction of new top-level DNS labels like me and search



Exercise Try looking up

- AAAA for bath.ac.uk
- AAAA for facebook.com
- SOA for bath.ac.uk
- A for moodle.bath.ac.uk
- And so on



Exercise And

• MX for bath.ac.uk

What is the Uni's mail server physical location?

What are the privacy/security/legal aspects?

Then try MX for bath.edu, an address the University uses for alumni (graduates and past staff)



Exercise Read about the public DNS servers like 1.1.1.1, 8.8.8.8 and others; why would you want to use them over your local DNS server?

Exercise What are the privacy/security/legal aspects of using public servers?

Exercise Find out how to buy a DNS name

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Now, UDP is fast but unreliable, and 512 byte datagrams won't be fragmented (recall: must be able to send 576 bytes IP), so there is little complication if a DNS datagram is lost: the source will just send a new request



UDP is preferred as it is fast with little overhead; while a TCP connection has a considerable setup cost



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So small and fast wherever possible; but slower and reliable if needed



DNS is good, but has problems

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- The spoof page could be made to look identical to the real bank web page, inviting me to enter my id and password



Exercise The security company RSA was attacked by DNS spoofing. Read about this

Exercise April 2018: routes to Amazon DNS servers were faked so that DNS requests were sent to fake servers. Read about this

Exercise Your home connection probably uses your ISP as a DNS server. ISPs have been known to rewrite DNS replies. They also might block access to other public DNS servers. Read about this



A partial solution exists in *Secure DNS* (DNSSec), which uses cryptography to authenticate DNS lookups



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Not much in use as it was introduced when people still thought that cryptography was too expensive to use



And neither DNS nor DNSSec provide any privacy: anyone listening can see what hosts you are trying to contact by monitoring your DNS requests



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This has a fair overhead over plain DNS, but provides both authentication and privacy

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But:

- The overhead can be managed fairly well (see HTTP *Keep-Alive* and TLS reconnections)
- While an uncooperative ISP could block the DNS over TLS assigned port (853) the HTTPS port (443) cannot be blocked without a lot of collateral damage to normal browsing



Exercise CIDR makes PTR lookups harder as netmasks, and therefore the delegations, are no longer on a byte boundary. Read RFC2317 on how CNAMES are used to solve this problem

Exercise Read about how (or if) DoH is supported in your browser

Exercise Read about DNS over Twitter and DNS over SMS