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Human factors are very important: we mustn't forget social engineering attacks

Also, there is no point in having a military-grade encryption system if you have an easily-guessable password

There are actually two problems to address:

- Secrecy
- Authentication

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Secrecy is familiar, but authentication is more fundamental

Secrecy is

make sure that this data is not readable by anyone other than the recipient

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Authentication is

make sure the recipient is who I think they are

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Perhaps you are buying over the Web and you want to send your credit card details to Acme Widget Company

So you negotiate a military-grade encryption key with them and send the details, happy in the knowledge that no-one else can read them

There is no point sending a strongly encrypted message only to find you sent it to the wrong person

Perhaps you are buying over the Web and you want to send your credit card details to Acme Widget Company

So you negotiate a military-grade encryption key with them and send the details, happy in the knowledge that no-one else can read them

Later you discover the Web page was a fake and your details are now in the hands of criminals

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So you negotiate a military-grade encryption key with them and send the details, happy in the knowledge that no-one else can read them

Later you discover the Web page was a fake and your details are now in the hands of criminals

You must have some way of determining if someone is *who they say they are*: this is authentication

In the real world we use documents like passports and driving licences to identify people

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In the Internet world we do the same, except now the documents are chunks of mathematical data

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In the Internet world we do the same, except now the documents are chunks of mathematical data

For details go to a crypto course!



If we were doing things properly, we should also talk about *authorisation* at this point



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After authentication there is the question of whether this entity is allowed access to some resource

Aside

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After authentication there is the question of whether this entity is allowed access to some resource

For example, in WPA-PSK, a correct password is usually taken as authorisation; in WPA-Enterprise the server will have a list of allowed users+passwords

We have already considered a link layer security: WPA

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We now look at a few others, including IPSec (network layer), PPTP and L2TP (link layer), SSL/TLS (transport layer)

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We now look at a few others, including IPSec (network layer), PPTP and L2TP (link layer), SSL/TLS (transport layer)

In the normal way of blurring layers when thinking about functionality, PPTP and L2TP are regarded as link layer, even though they layer over IP

IPSec can be used to set up secure point-to-point links (see VPNs, later), but can also be used to secure and authenticate individual connections when the other end supports it: *opportunistic encryption*

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IPSec consists of several protocols and defines several IP optional header fields

Secrecy is implemented by *Encapsulating Security Payload* (ESP)

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Authentication by Authentication Header (AH)

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Authentication by Authentication Header (AH)

Keys are managed by Internet Key Exchange (IKE) which itself uses the Internet Security Association and Key Management Protocol (ISAKMP)

IPSec authenticates connections, not users

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You do not use it to login, but to ensure the remote host really is Acme Widgets before you send data (money) to them

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Both ESP and AH require a secret shared key to work

This key can be

- pre-agreed (manual keying)
- negotiated by IKE

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IKE can itself use a pre-agreed key to deliver the ESP/AH key, or use a public-key certificate mechanism

Normally there is one IKE process per host and it manages all exchanges for that host

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When a new IPSec/IP connection is started an IKE exchange will take place before the IPSec can continue

This may take some time: even enough to cause a TCP timeout on slow machines

IPSec

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- is directly inside the IP layer (optional headers), so UDP and TCP are easily layered transparently on top
- clearly only applies to IP
- AH does authentication, while ESP does secrecy and authentication. Pure authentication is OK if you do not need secrecy, but pure secrecy is open to impersonation attacks without some authentication

• ESP has a trailer as well as a header: this can contain padding to hide the length of the original packet

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- ESP only authenticates the payload, while AH authenticates all the packet, excepting the mutable fields (TTL etc.) that change en route

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- ESP only authenticates the payload, while AH authenticates all the packet, excepting the mutable fields (TTL etc.) that change en route
- applies to both IPv4 and IPv6

Problems:

• initial connection overhead is high

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- IPSec is tricky to set up and manage
- It works at the OS host level, and so needs a competent administrator
- and also does not have the flexibility that (say) SSL/TLS has, allowing each application to be managed independently

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But the configuration question remains

Exercise The University uses IPsec: investigate

A new alternative to IPSec that is growing in popularity is *WireGuard*

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Exercise Read about WireGuard

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Some systems are based around creating a *Virtual Private Networks* (VPN)

A VPN allows a machine to appear to be on another network by means of tunnelling

Recall tunnelling: where one protocol is layered over another so the lower protocol can transport the upper protocol transparently over a network that might not normally carry the upper protocol

VPNs are *private* as they add encryption of the data in the tunnel to provide security

Traffic from the host travels through the tunnel to the network, where it can be routed as if it had originated there

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This allows the host to use the services on the network as if it were local to that network

This is good for *teleworkers*

For example, IPSec and WireGuard are VPNs

In overview a VPN:

• is software that creates a new virtual network interface

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- when packets need to use the VPN they are routed out on that interface
- the packets are encrypted (in the kernel or in user level software)
- this data is now sent out over a real interface, e.g., using UDP

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As far as the client and server are concerned, they are operating as normal, sending and receiving packing on a normal interface

There are two common setups for VPNs that treat data from the client in different ways

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• where *all* client traffic goes out through the VPN, and the server end of the VPN routes it onwards as appropriate

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- where *all* client traffic goes out through the VPN, and the server end of the VPN routes it onwards as appropriate
- where only traffic destined for the server's network goes through the VPN. Other traffic from the source goes directly to its destination in the normal way. This is sometimes called a *split tunnel*

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
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fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0
fire.cs.bath.ac				0	0	

A packet destined for the CS network goes through (virtual) interface tun0, which actually sends packets to the VPN software on the local machine;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
• • •						
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0

This is encrypted and encapsulated in a packet with destination fire on the CS network;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0

And port number that of the VPN software on the remote server;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
• • •						
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0
••••						

This packet then goes through the host (H) route to CS through the real interface wlan0;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0

The routes are checked longest mask first, to prevent an infinite loop!

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
• • •						
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0
••••						

The VPN software on the remote server gets the packet and deencapsulates it;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
• • •						
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0
fire.cs.bath.ac				0	0	

It rewrites the source address on the packet to its own address so that replies come back to it (c.f., NAT);

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
• • •						
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0
••••						

The remote host forwards the packet to the destination which is on its local network;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0

Symmetrically, it will encrypt, encapsulate and return any replies back through the tunnel;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
• • •						
default	home.gateway.ho	0.0.0.0	UG	0	0	0 wlan0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0
fire.cs.bath.ac				0	0	

Locally, all other (default) traffic goes out through the normal interface, not the VPN

Destination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface
default	172.16.0.1	0.0.0.0	UG	0	0	0	tun0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0	tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0	wlan0

In comparison, the other setup;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface
default	172.16.0.1	0.0.0.0	UG	0	0	0	tun0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0	tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0	wlan0
•••							

Here all default traffic goes through interface tun0;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface
• • •							
default	172.16.0.1	0.0.0.0	UG	0	0	0	tun0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0	tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0	wlan0
••••							

This is encapsulated in the same way;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface
default	172.16.0.1	0.0.0.0	UG	0	0	0	tun0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0	tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0	wlan0

And routed through the tunnel via the real interface;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use Iface
default	172.16.0.1	0.0.0.0	UG	0	0	0 tun0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0 tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0 wlan0
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The VPN software on the remote network gets the packet, deencapsulates it, rewrites its source address and forwards it to the destination which now might be anywhere, not just on its local network;

Destination	Gateway	Genmask	Flags	Metric	Ref	Use	Iface
default	172.16.0.1	0.0.0.0	UG	0	0	0	tun0
cs.bath.ac.uk	172.16.0.1	255.255.255.240	UG	2	0	0	tun0
fire.cs.bath.ac	home.gateway.ho	255.255.255.255	UGH	0	0	0	wlan0

Returning packets are sent back through the VPN as before

In the first setup you get secure access to the work network, but to the rest of the world you are at home (and not secure)

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The second makes you to appear to everyone to be at work as all your packets have a work IP address on them

Exercise Find out how to use the University's VPN from your home

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Exercise The www.newscientist.com website reserves some content for subscribers. It uses the requesting IP address to check for access rights. The University is a subscriber. Use a VPN to get at this content from your home

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Exercise The www.newscientist.com website reserves some content for subscribers. It uses the requesting IP address to check for access rights. The University is a subscriber. Use a VPN to get at this content from your home

Exercise Find out how (or if) your favourite VPN can be configured as a split tunnel

The *Point-to-Point Tunneling Protocol* (PPTP) was devised by Microsoft to support VPNs

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- tunnels IP over PPP over the *Generic Routing Encapsulation* protocol (GRE) over IP and sends connection control messages over a separate TCP connection
- layers only over IP

lt

 can encapsulate other protocols such as IPX (Internetwork Packet Exchange, Novell) and NetBIOS/NetBEUI (Network BIOS, NetBIOS Extended User Interface, Microsoft)

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- uses PPP for authentication
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- is simple to set up

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A later version (MS-CHAPv2) fixes some, but not all of the holes

The *Layer 2 Tunneling Protocol* (L2TP) combines features of PPTP and *Layer 2 Forwarding* (L2F) developed by Cisco Systems Inc

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lt

tunnels IP over PPP over L2TP over UDP

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- is intended to be used over ATM, Frame Relay and X.25 networks

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We shall talk about OpenVPN, later

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Plus, each TCP has its own idea of timeouts and retransmits and they can start to fight each other: the retransmit of one TCP will be viewed as a duplicate by the other TCP

Thus most VPNs tunnel over UDP

Exercise Read about the *Secure Socket Tunneling Protocol* (SSTP) and the *TCP meltdown problem*

Generic problems of VPNs include

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- there is encryption and authentication header overhead in every packet: this may cause extra packets or extra fragmentation
- there is overhead in the time taken to encrypt or authenticate the packets

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- some ISPs like to charge more for, or manage certain kinds of traffic (e.g., bittorrent and video) and this hides the kind of traffic. So some ISPs have blanket bans on VPNs
- in VPNs speed is secondary to security, but people will not use them if they are too slow

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Note that a VPN can make you appear to be in a different country

Good for evading country-locked content (geofencing), but bad for law enforcement and people who want to track what you are doing

For example, the *Investigatory Powers Act 2016*, and its update, *The Data Retention and Acquisition Regulations 2018*, a law that can require your ISP to log every website you visit and every recipient of emails and phone calls (your *Internet Connection Records*)

Some core information is accessible, without warrant, by certain people:

- account reference
- a source and port address, a destination IP and port address
- a time/date + duration
- partial URLs (only part containing server name, no content)

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- account reference
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- a time/date + duration
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So, the metadata and not the data

An interception warrant is needed for more, e.g., content

The law also gives the security services new powers to hack computers and, e.g., pressure service providers not to support end-to-end encryption

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Exercise Read the list of 50 or so authorities that can access your web history, without warrant